

Implementing preferences with *asprin*

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Outline

- 1 Introduction
- 2 Preliminaries
- 3 Language
- 4 Implementation
- 5 Embedding existing approaches
- 6 Experimental analysis
 - Boosting optimization via heuristics
 - Dedicated systems versus *asprin*
- 7 Summary

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Motivation

- The identification of preferred, or optimal, solutions is often indispensable in real-world applications

In many cases, this also involves the combination of various qualitative and quantitative preferences

- Only optimization statements representing objective functions using summation are established components of today's ASP systems
- Example `#minimize{40 : sauna, 70 : dive}`

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- Example $\#minimize\{40 : sauna, 70 : dive\}$

Approach

- asprin is a framework for handling preferences among the (stable) models of logic programs
 - general because it captures many existing approaches to preference
 - flexible because it allows an easy implementation of new approaches
- asprin builds upon advanced control capacities for multi-shot and meta solving, allowing for
 - reducing redundancies via ordinary ASP encodings

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Example

Your vacation (logic) program ...

... talking about *sauna*, *dive*, *hike*, *bunji*, *hot*, etc

```
#preference(costs, less(weight)){40 : sauna, 70 : dive}
#preference(fun, superset){sauna, dive, hike, ¬bunji}
#preference(temps, aso){dive > sauna || hot, sauna > dive || ¬hot}
#preference(all, pareto){name(costs), name(fun), name(temps)}
#optimize(all)
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Preference

- A strict partial order $\succ$ on the stable models of a logic program
That is, $X \succ Y$ means that X is preferred to Y
- A stable model X is $\succ$ -preferred, if there is no other stable model Y such that $Y \succ X$
- A preference type is a (parametric) class of preference relations
- Example *subset*, *pareto*, etc.

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Language

- naming atom *name(s)*
 where s is the name of a preference
- weighted formula $w_1, \dots, w_l : \phi$
 where each w_i is a term and ϕ is a Boolean formula
- preference element $\Phi_1 > \dots > \Phi_m \parallel \Phi$
 where each Φ_r is a set of weighted formulas and Φ is a non-weighted formula
- preference statement $\# \text{preference}(s, t) \{e_1, \dots, e_n\}$
 where s and t represent the preference statement and its type
 and each e_j is a preference element
- optimization directive $\# \text{optimize}(s)$
 where s is the name of a preference
- preference specification is a set S of preference statements and a directive
 $\# \text{optimize}(s)$ such that S is acyclic, closed, and $s \in S$

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Preference type

- A **preference type** t is a function mapping a set of preference elements E to a preference relation

$$t(E) \subseteq \mathcal{A} \times \mathcal{A}$$

- Examples

- $(X, Y) \in \text{subset}(E)$ iff $\{\ell \in E \mid X \models \ell\} \subset \{\ell \in E \mid Y \models \ell\}$
 - $(X, Y) \in \text{pareto}(E)$ iff $\bigwedge_{name(s) \in E} (X \succeq_s Y) \wedge \bigvee_{name(s) \in E} (X \succ_s Y)$

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Preference relation

- A **preference relation** is obtained by applying a preference type to a set of preference elements

$\#preference(s, t) E$ declares preference relation $t(E)$, denoted by $\succ_s$

- Example $\#preference(1, subset)\{a, b, c\}$ declares

$$X \succ_1 Y \text{ iff } \{\ell \in \{a, b, c\} \mid X \models \ell\} \subset \{\ell \in \{a, b, c\} \mid Y \models \ell\}$$

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Preference program

- Reification $H_X = \{holds(a) \mid a \in X\}$ and $H'_X = \{holds'(a) \mid a \in X\}$
- Preference program Let s be a preference statement declaring $\succ_s$.

We define Q_s as a preference program for s , if for all sets $X, Y \subseteq \mathcal{A}$, we have

$$X \succ_s Y \text{ iff } Q_s \cup H_X \cup H'_Y \text{ is satisfiable}$$

- Note Q_s is implemented as $F_s \cup E_{t_s} \cup C$
- Note *asprin's* expressiveness is delineated by the decision problem

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*#*preference(1, subset){a, b, c}
*#*optimize(1)

$$H_{\{a\}} = \left\{ \text{holds}(a). \right\}$$

$$H'_{\{a,b\}} = \left\{ \text{holds}'(a). \quad \text{holds}'(b). \right\}$$

$$F_1 = \left\{ \begin{array}{ll} \text{preference}(1, \text{subset}). & \text{preference}(1, 1, 1, \text{for}(a), ()). \\ \text{optimize}(1). & \text{preference}(1, 2, 1, \text{for}(b), ()). \\ & \text{preference}(1, 3, 1, \text{for}(c), ()). \end{array} \right\}$$

$$E_{\text{subset}} = \left\{ \begin{array}{l} \text{better}(P) :- \text{preference}(P, \text{subset}), \\ \quad \text{not holds}(X), \quad \text{holds}'(X), \text{preference}(P, _, _, \text{for}(X), _), \\ \quad \text{not holds}(Y) : \text{not holds}'(Y), \text{preference}(P, _, _, \text{for}(Y), _). \end{array} \right\}$$

$$C = \left\{ :- \text{optimize}(P), \text{not better}(P). \right\}$$

There is a stable model, indicating that $\{a\} \succ_1 \{a, b\}$

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There is no stable model, indicating that $\{a, b\} \not\preceq_1 \{a\}$

Basic algorithm $solveOpt(P, s)$

Input : A program P over $\mathcal{A}$ and preference statement s

Output : A $\succ_s$ -preferred stable model of P , if P is satisfiable, and $\perp$ otherwise

$Y \leftarrow solve(P)$

if $Y = \perp$ **then return** $\perp$

repeat

$X \leftarrow Y$
 $Y \leftarrow solve(P \cup Q_s \cup R \cup H'_X)$

until $Y = \perp$

return X

where $R = \{holds(a) \leftarrow a \mid a \in \mathcal{A}\}$

asprin's library

■ Basic preference types

- `subset` and `superset`
- `less(cardinality)` and `more(cardinality)`
- `less(weight)` and `more(weight)`
- `aso` (Answer Set Optimization)
- `poset` (Qualitative Preferences)

■ Composite preference types

- `neg`
- `and`
- `pareto`
- `lexico`

■ Customized preference types

 See paper (and Potassco Guide) on how to define new types!

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Available in *asprin*, continued

- Weak constraints (DLV) and minimize statements (SMODELS)
- Answer set optimization (Brewka, Niemelä & Truszczyński; IJCAI)
 - Answer set optimization with penalties (Brewka; KR)
 - Ordered disjunctions (Brewka; AAI)
- Qualitative preferences (Di Rosa, Giunchiglia & Maratea; Constraints)
- Preference language (Son & Pontelli; TPLP)

 *See paper for details and experiments next!*

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Heuristic framework

- *clingo* allows for incorporating domain-specific heuristics into ASP solving
 - input language for expressing domain-specific heuristics
 - solving capacities for integrating domain-specific heuristics
- Heuristic predicate `_heuristic`
- Heuristic modifiers (atom, *a*, and integer, *v*)
 - `init` for initializing the heuristic value of *a* with *v*
 - `factor` for amplifying the heuristic value of *a* by factor *v*
 - `level` for ranking all atoms; the rank of *a* is *v*
 - `sign` for attributing the sign of *v* as truth value to *a*
- Heuristic atoms
 - `_heuristic(occurs(A,T),factor,T) :- action(A), time(T).`

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- Heuristic predicate `_heuristic`
- Heuristic modifiers (atom,  $a$ , and integer,  $v$ )
  - `init` for initializing the heuristic value of  $a$  with  $v$
  - `factor` for amplifying the heuristic value of  $a$  by factor  $v$
  - `level` for ranking all atoms; the rank of  $a$  is  $v$
  - `sign` for attributing the sign of  $v$  as truth value to  $a$
- Heuristic atoms
  - `_heuristic(occurs(A,T),factor,T) :- action(A), time(T).`

# Heuristic framework

- *clingo* allows for incorporating domain-specific heuristics into ASP solving
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  - Heuristic atoms
- ```
_heuristic(occurs(mv,5),factor,5) :- action(mv), time(5).
```

asprin with different heuristic settings

Benchmark \ System	<i>asprin_w</i>	<i>asprin_w+s</i>	<i>asprin_w+l</i>	<i>asprin_w+f</i>	<i>asprin_s</i>	<i>asprin_s+s</i>	<i>asprin_s+l</i>	<i>asprin_s+f</i>
<i>Ricochet</i> (30) 20.00	432 (8, 4)	407 (7, 4)	68 (1, 0)	71 (1, 0)	365 (8, 3)	461 (7, 10)	69 (1, 0)	71 (1, 0)
<i>Timetabling</i> (12) 23687.75	345 (285, 3)	255 (202, 2)	900 (4, 12)	6 (1, 0)	217 (144, 2)	21 (18, 0)	900 (2, 12)	5 (1, 0)
<i>Puzzle</i> (7) 580.57	82 (2, 0)	112 (2, 0)	136 (2, 0)	416 (2, 1)	31 (1, 0)	32 (1, 0)	21 (1, 0)	51 (1, 0)
<i>Crossing</i> (24) 211.92	104 (42, 1)	98 (35, 0)	805 (19, 20)	387 (6, 6)	0 (6, 0)	1 (6, 0)	7 (9, 0)	3 (1, 0)
<i>Valves</i> (30) 56.63	69 (7, 0)	65 (6, 0)	460 (8, 11)	715 (0, 22)	38 (4, 0)	39 (4, 0)	339 (4, 6)	673 (0, 21)
<i>Expansion</i> (30) 7501.87	216 (299, 0)	10 (15, 0)	38 (7, 0)	12 (3, 0)	64 (295, 0)	14 (54, 0)	4 (4, 0)	3 (1, 0)
<i>Repair</i> (30) 6750.73	76 (48, 0)	15 (47, 0)	71 (3, 2)	8 (2, 0)	8 (43, 0)	3 (11, 0)	1 (1, 0)	1 (1, 0)
<i>Diagnosis</i> (30) 1669.00	196 (341, 3)	76 (66, 0)	43 (4, 0)	118 (3, 2)	19 (338, 0)	2 (39, 0)	0 (1, 0)	0 (1, 0)
$\emptyset(\emptyset, \Sigma)$	190 (129, 11)	130 (48, 6)	315 (6, 45)	217 (2, 31)	93 (105, 5)	72 (18, 10)	168 (3, 18)	101 (1, 21)

■ *asprin_w*■ *asprin_w+s*■ *asprin_w+l*■ *asprin_w+f*

■ w — weight-based

■ s — sign heuristic

■ l — level-based heuristic

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asprin with different heuristic settings

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<i>Ricochet</i> (30)	20.00	432 (8, 4)	407 (7, 4)	68 (1, 0)	71 (1, 0)	365 (8, 3)	461 (7, 10)	69 (1, 0)	71 (1, 0)
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asprin with different heuristic settings

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Outline

- 1 Introduction
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 - **Dedicated systems versus asprin**
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asprin versus *clingo* and *metasp*

(B,D,R&S; AAAI)

Benchmark \ System	<i>clingo</i>	<i>asprin_w</i>	<i>asprin_w -f</i>	<i>metasp</i>	<i>asprin_s</i>	<i>asprin_s -f</i>
<i>Ricochet</i> (30) 20.00	104.74 (0)	174.26 (0)	113.45 (0)	811.32 (24)	175.71 (0)	109.91 (0)
<i>Timetabling</i> (12) 23687.75	35.82 (0)	490.39 (5)	694.92 (8)	798.75 (10)	142.03 (0)	12.01 (0)
<i>Puzzle</i> (7) 580.57	77.00 (0)	77.39 (0)	96.70 (0)	34.79 (0)	17.06 (0)	17.22 (0)
<i>Crossing</i> (24) 211.92	48.43 (0)	105.64 (1)	67.50 (0)	62.33 (0)	0.50 (0)	0.46 (0)
<i>Valves</i> (30) 56.63	52.53 (0)	72.85 (0)	78.11 (0)	900.00 (30)	45.01 (0)	39.31 (0)
<i>Expansion</i> (30) 7501.87	91.53 (0)	373.56 (2)	241.05 (7)	900.00 (30)	292.57 (0)	21.12 (0)
<i>Repair</i> (30) 6750.73	71.78 (0)	102.19 (0)	43.94 (0)	900.00 (30)	6.88 (0)	2.19 (0)
<i>Diagnosis</i> (30) 1669.00	84.96 (0)	254.19 (3)	101.33 (0)	181.71 (6)	41.55 (0)	1.56 (0)
$\emptyset(\Sigma)$	70.85 (0)	206.31 (11)	179.63 (15)	573.61 (130)	90.16 (0)	25.47 (0)

■ *clingo* (using branch-and-bound)■ *asprin_w*■ *asprin_w -f*

■ w — weight-based

■ -f — no phase saving

■ *metasp* (using disjunction)■ *asprin_s*■ *asprin_s -f*

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asprin versus clingo and metasp

(B,D,R&S; AAAI)

Benchmark \ System	<i>clingo</i>	<i>asprin_w</i>	<i>asprin_w -f</i>	<i>metasp</i>	<i>asprin_s</i>	<i>asprin_s -f</i>
<i>Ricochet</i> (30) 20.00	104.74 (0)	174.26 (0)	113.45 (0)	811.32 (24)	175.71 (0)	109.91 (0)
<i>Timetabling</i> (12) 23687.75	35.82 (0)	490.39 (5)	694.92 (8)	798.75 (10)	142.03 (0)	12.01 (0)
<i>Puzzle</i> (7) 580.57	77.00 (0)	77.39 (0)	96.70 (0)	34.79 (0)	17.06 (0)	17.22 (0)
<i>Crossing</i> (24) 211.92	48.43 (0)	105.64 (1)	67.50 (0)	62.33 (0)	0.50 (0)	0.46 (0)
<i>Valves</i> (30) 56.63	52.53 (0)	72.85 (0)	78.11 (0)	900.00 (30)	45.01 (0)	39.31 (0)
<i>Expansion</i> (30) 7501.87	91.53 (0)	373.56 (2)	241.05 (7)	900.00 (30)	292.57 (0)	21.12 (0)
<i>Repair</i> (30) 6750.73	71.78 (0)	102.19 (0)	43.94 (0)	900.00 (30)	6.88 (0)	2.19 (0)
<i>Diagnosis</i> (30) 1669.00	84.96 (0)	254.19 (3)	101.33 (0)	181.71 (6)	41.55 (0)	1.56 (0)
$\emptyset(\Sigma)$	70.85 (0)	206.31 (11)	179.63 (15)	573.61 (130)	90.16 (0)	25.47 (0)

■ *clingo* (using branch-and-bound)

■ *asprin_w*

■ *asprin_w -f*

■ *w* — weight-based

■ *-f* — no phase saving

■ *metasp* (using disjunction)

■ *asprin_s*

■ *asprin_s -f*

■ *s* — subset-based

■ *-f* — no phase saving

aso versus *asprin*

n	<i>aso</i>	<i>aso_l</i>	<i>asprin_a</i>	<i>asprin_{l+a}</i>
350	9 (0)	17 (0)	4 (0)	5 (0)
360	14 (0)	22 (0)	48 (0)	50 (0)
370	15 (0)	25 (0)	38 (0)	39 (0)
380	10 (0)	23 (0)	8 (0)	9 (0)
390	59 (0)	72 (0)	50 (1)	52 (1)
400	22 (0)	33 (0)	28 (0)	30 (0)
410	87 (1)	96 (1)	124 (2)	125 (2)
420	97 (1)	108 (1)	60 (0)	62 (0)
430	68 (0)	79 (0)	144 (0)	147 (0)
440	165 (3)	175 (3)	165 (2)	167 (2)
450	45 (0)	61 (0)	52 (0)	54 (0)
460	112 (1)	125 (1)	117 (2)	120 (2)
470	201 (4)	210 (4)	161 (2)	162 (2)
480	152 (2)	165 (2)	70 (1)	72 (1)
490	206 (2)	218 (2)	265 (4)	267 (4)
$\emptyset(\Sigma)$	84 (14)	95 (14)	89 (14)	91 (14)

■ *aso* — dedicated system

■ *asprin_a*

■ *aso_l* — dedicated system
with level-based heuristics

■ *asprin_{l+a}* — with level-based heuristics

satpref versus *asprin*

Benchmark \ System	<i>satpref</i>	<i>satpref</i> +s	<i>satpref</i> +H	<i>asprin_p</i>	<i>asprin_p</i> +s	<i>asprin_p</i> +H
0.0	0 (29, 0)	0 (1, 0)	0 (1, 0)	1 (16, 0)	0 (2, 0)	0 (1, 0)
0.00621	0 (35, 0)	0 (1, 0)	90 (1, 6)	1 (17, 0)	1 (2, 0)	1 (1, 0)
0.01243	1 (75, 0)	1 (3, 0)	118 (1, 7)	6 (26, 0)	2 (3, 0)	3 (1, 0)
0.02486	8 (388, 0)	6 (10, 0)	635 (1, 38)	55 (74, 0)	9 (8, 0)	64 (1, 4)
0.04972	67 (1463, 2)	16 (36, 0)	900 (0,100)	318 (203, 16)	26 (17, 0)	176 (1,14)
1.0	850 (10315,88)	243 (590,10)	177 (1, 12)	856 (323, 92)	174 (96, 0)	280 (1,24)
$\emptyset(\emptyset, \Sigma)$	154 (2051,90)	44 (107,10)	320 (1,163)	206 (110,108)	35 (21, 0)	88 (1,42)
MAXSAT	54 (8849, 0)	9 (7, 0)	62 (1, 0)	835 (957, 31)	109 (31, 3)	171 (1, 6)
PBO/ <i>pbo-mqc-nencdr</i>	5 (267, 0)	2 (2, 0)	664 (1, 88)	150 (207, 14)	9 (2, 0)	244 (1,20)
PBO/ <i>pbo-mqc-nlogencdr</i>	3 (228, 0)	1 (2, 0)	237 (1, 21)	110 (214, 3)	5 (2, 0)	141 (1,15)
PSEUDO/ <i>primes</i>	110 (396,18)	110 (1,18)	110 (1, 18)	215 (334, 27)	106 (5,17)	110 (1,17)
PSEUDO/ <i>routing</i>	346 (409, 4)	49 (1, 0)	50 (1, 0)	85 (475, 0)	4 (1, 0)	86 (1, 1)
Partial-MINONE	14 (2, 0)	14 (2, 0)	7 (1, 0)	24 (2, 0)	4 (1, 0)	25 (1, 0)
$\emptyset(\emptyset, \Sigma)$	88 (1692,22)	31 (2,18)	188 (1,127)	236 (365, 75)	43 (7,20)	129 (1,59)

■ *satpref*■ *satpref*+s■ *satpref*+H

■ s — sign heuristic

■ H — complex heuristic

■ *asprin_p*■ *asprin_p*+s■ *asprin_p*+H

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Summary

- *asprin* stands for “ASP for Preference handling”
- *asprin* is a general, flexible, and extendable framework for preference handling in ASP
- *asprin* caters to
 - off-the-shelf users using the preference relations in *asprin*'s library
 - preference engineers customizing their own preference relations
- *asprin* can be boosted by *clingo*'s heuristic framework
- <http://potassco.sourceforge.net/labs.html#asprin>

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